

Randomness for Ergodic Measures

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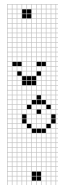
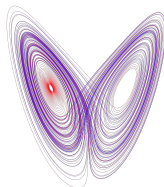
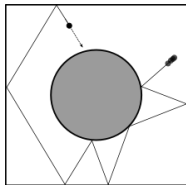
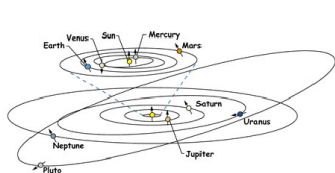
Slides available at
www.personal.psu.edu/jmr71/

(Updated on July 3, 2015.)

Ergodic theory

Dynamical systems

- Dynamical systems is the study of “motion”.



- Concerns a **space** equipped with an **action**.
 - Topological space with continuous action vs. **probability space with a measure-preserving action**.
 - Continuous time dynamics vs. **discrete time dynamics**.

Symbolic dynamics

- **Space:** Cantor space $\{0,1\}^{\mathbb{N}}$ (space of infinite binary sequences)
- **Transformation:** Left shift map $S: \{0,1\}^{\mathbb{N}} \rightarrow \{0,1\}^{\mathbb{N}}$

$$0x, 1x \mapsto x.$$

- **Measure:** A Borel probability measure μ on $\{0,1\}^{\mathbb{N}}$ given by

$$\begin{aligned} \mu(\sigma) &= \text{probability that a sequence starts with } \sigma \quad (\sigma \in \{0,1\}^*) \\ \mu(\sigma) &\geq 0, \quad \mu(\sigma 0) + \mu(\sigma 1) = \mu(\sigma), \quad \mu(\text{ }) = 1. \end{aligned}$$

- A measure is computable if $\mu(\sigma)$ is computable from σ .
- We also require that the measure be invariant under the shift action.

Shift-invariant measures

- The fair-coin measure $\lambda(\sigma) = (1/2)^{|\sigma|}$ satisfies the following:
If one removes the first bit, one still has the same probability measure.
- Such a measure is called **shift-invariant**.
- More formally, μ is shift invariant if
 - $\mu(S^{-1}A) = \mu(A)$ for all measurable $A \subseteq \{0,1\}^{\mathbb{N}}$ (for left shift map S).
 - Equivalently, $\mu(0\sigma) + \mu(1\sigma) = \mu(\sigma)$ for all finite strings $\sigma \in \{0,1\}^*$.
- There are a wide and complex variety of shift invariant measures.
 - Bernoulli processes (i.i.d. weighted coin flips)
 - Markov processes (under certain initial conditions)
 - and many many more

Combining shift-invariant measures

- If μ_1 and μ_2 are shift-invariant measures, then the convex combination

$$\frac{1}{3}\mu_1 + \frac{2}{3}\mu_2$$

is shift-invariant.

- The same is true of any convex combination of shift-invariant measures.
- However, the fair-coin measure satisfies this property.

The fair-coin measure is not a (nontrivial) convex combination of two or more different shift-invariant measures.

- A shift-invariant measure that cannot be decomposed is called a **ergodic measure**.
- Also, μ is ergodic iff for all sets A , if $T^{-1}(A) = A$ then $\mu(A)$ is 0 or 1.

Intermission: Algorithmic randomness

Algorithmic randomness

- Informally, a point is random if it satisfies every “effective probability one property”.
- For a *computable* probability measure μ on $\{0,1\}^{\mathbb{N}}$:
 - A **μ -Martin-Löf test** is a sequence (U_n) of effectively open (Σ_1^0) sets in $\{0,1\}^{\mathbb{N}}$ such that $\mu(U_n) \leq 2^{-n}$.
 - A **μ -Martin-Löf random** is a point not in $\bigcap_n U_n$ for any Martin-Löf tests (U_n) .
- For a *non-computable* probability measure μ on $\{0,1\}^{\mathbb{N}}$:
 - A **μ -Martin-Löf random** is a point not in $\bigcap_n U_n$ for any Martin-Löf tests (U_n) *relativized* to the measure μ .
 - A **blind- μ -Martin-Löf random** is a point not in $\bigcap_n U_n$ for any Martin-Löf tests (U_n) (*not relativized*).

Back to our program: Decomposing measures

Substring frequencies / empirical ergodic measures

- The frequency of times that the string σ appears in the sequence x is

$$\text{freq}(\sigma, x) = \lim_n \frac{\#\{k \leq n : x \text{ contains substring } \sigma \text{ at position } k\}}{n}.$$

- E.g., x is normal iff $\text{freq}(\sigma, x) = (1/2)^{|\sigma|} = \lambda(\sigma)$ (fair-coin measure).
- If μ is shift-invariant,
 - $\text{freq}(\sigma, x)$ exists for μ -almost-every x (pointwise ergodic theorem).
 - $\text{freq}(\cdot, x)$ is μ -almost-surely an ergodic measure.
 - If μ is ergodic, then μ -almost-surely $\text{freq}(\sigma, x) = \mu(\sigma)$.
 - $\text{freq}(\cdot, x)$ is also called the **empirical ergodic measure**.
 - μ is a unique convex combination of ergodic measures:

$$\mu(\sigma) = \int \text{freq}(\sigma, x) d\mu(x).$$

Bucket of weighted-coins

- Bucket of weighted-coins process:
 - Start with a bucket of weighted coins (under some distribution).
 - Step -1: Take a coin at random from the bucket.
 - Steps 0,1,2,...: Flip coin and record outcome. (Same coin every time.)
- The resulting measure μ is
 - shift-invariant, hence $\mu(\sigma) = \int \text{freq}(\sigma, x) d\mu(x)$.
 - **exchangeable**, i.e. invariant under finite permutations of bits.
 - $\text{freq}(\cdot, x)$ is μ -almost-surely a Bernoulli (weighted coin) measure.
 - μ is a convex combination of Bernoulli measures
- de Finetti's theorem: Every exchangeable measure comes from "a bucket of weighted coins".
- If μ is computable and exchangeable and x is μ -random,
 - $\text{freq}(\cdot, x)$ exists and is a Bernoulli measure.
 - $\text{freq}(\cdot, x)$ is computable from x .
 - x is $\text{freq}(\cdot, x)$ -random.
 - For each Bernoulli measure ν , blind- ν -randomness = ν -randomness.

Main Questions and Answers

- We have the analogy

$$\frac{\text{ergodic}}{\text{shift-invariant}} = \frac{\text{Bernoulli}}{\text{exchangable}}.$$

- Assume μ is computable and shift-invariant and x is μ -random.
 - Does $\text{freq}(\cdot, x)$ exist, and is it a measure?
 - Yes and yes. (V'yugin)
 - Is $\text{freq}(\cdot, x)$ ergodic?
 - **Still open.**
 - Is $\text{freq}(\cdot, x)$ computable from x ?
 - No. (Hoyrup)
 - Is x $\text{freq}(\cdot, x)$ -random (or ν -random for some ergodic ν)?
 - **No (and No). (Reimann and R.) (also Monin).**
 - For ergodic ν , does ν -random = blind- ν -random.
 - **No. (Reimann and R.) (also Monin).**
- (The answers are all yes if the “ergodic decomposition is computable.”)

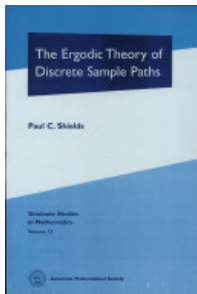
Proof Method (Cutting and Stacking)

Proof sketch

- Construct:
 - A computable shift-invariant measure μ
 - a μ -random x
- Such that:
 - $\text{freq}(\cdot, x)$ is an ergodic measure
 - x is blind- $\text{freq}(\cdot, x)$ -random
 - $\text{freq}(\cdot, x)$ computes x
- Hence:
 - x is μ -random and blind- $\text{freq}(\cdot, x)$ -random, but
 - x is not $\text{freq}(\cdot, x)$ -random (nor ν -random for any ergodic ν).
- Method:
 - Cutting and stacking

Cutting and stacking

- Originated in ergodic theory.
- Has been used by V'yugin and Franklin/Towsner in randomness.
- Well-suited to computability theorists.
- See pictures on the board.
- For more information see:
 - Paul C. Shields. *The Ergodic Theory of Discrete Sample Paths*.



Closing Thoughts

Thank You!

These slides will be available on my webpage:

<http://www.personal.psu.edu/jmr71/>

Or just Google™ me, “Jason Rute”.

P.S. I am on the job market.